

Type System for Egison

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Syntax

c denotes a character.

s denotes a string.

n denotes a number.

b denotes a boolean.

C denotes a constructor of an inductive data.

C_p denotes a constructor of an inductive pattern.

C_{dp} denotes a constructor of an inductive data pattern.

Expressions

$M, N ::= x \mid c \mid s \mid b \mid n$	variables and constants
$\mid (\text{if } M M M)$	
$\mid (\text{lambda } [\$x \$y \dots] M)$	abstraction
$\mid (M M \dots)$	application
$\mid [M \dots]$	tuple
$\mid \{M \dots\}$	collection
$\mid \langle C M \dots \rangle$	inductive data
$\mid \text{something}$	something matcher
$\mid (\text{match-all } M M [p M])$	
$\mid (\text{matcher } \{\phi \dots\})$	

Patterns

$p ::= -$	wildcard
$\mid \$x$	pattern variable
\mid ,M	value pattern
$\mid \langle C_p p \dots \rangle$	inductive pattern

Matcher Clauses

$\phi ::= [pp M \{[dp M]\}]$

Primitive Pattern Patterns

$pp ::= \$$	pattern hole
$\mid ,\$x$	value-pattern pattern
$\mid \langle C_{pp} pp \dots \rangle$	inductive-pattern pattern

Primitive Data Patterns

$dp ::= \$x$	primitive pattern variable
$\mid \langle C_{dp} dp \dots \rangle$	inductive-data pattern

Types

$S, T ::= X \mid I \mid \text{String} \mid \text{Integer} \mid \text{Bool}$
$\mid T \rightarrow T$
$\mid [T T \dots]$
$\mid \{T\}$
$\mid (\text{Pattern } T)$
$\mid (\text{PPattern } T T)$
$\mid (\text{PDPattern } T)$
$\mid (\text{Matcher } T)$

Contexts

$\Gamma ::= \epsilon$	empty context
$\mid \Gamma, x : T$	append x to the context
$\mid \Gamma ++ \Gamma$	join two contexts

Typing Rules

$$\begin{array}{c}
\frac{}{\Gamma, x : T, \Delta \vdash x : T} \text{T-VAR} \quad \frac{}{\Gamma \vdash s : \text{String}} \text{T-STR} \quad \frac{}{\Gamma \vdash n : \text{Integer}} \text{T-NUM} \quad \frac{}{\Gamma \vdash b : \text{Bool}} \text{T-BOOL} \\
\frac{\Gamma \vdash M_1 : \text{Bool} \quad \Gamma \vdash M_2 : T \quad \Gamma \vdash M_3 : T}{\Gamma \vdash (\text{if } M_1 M_2 M_3) : T} \text{T-IF} \\
\frac{\Gamma \vdash \{x_1 : S_1, x_2 : S_2, \dots, x_n : S_n\} \vdash M : T}{\Gamma \vdash (\text{lambda } [\$x_1 \$x_2 \dots \$x_n] M) : [S_1, S_2, \dots] \rightarrow T} \text{T-ABS} \\
\frac{\Gamma \vdash M : [S_1 S_2 \dots S_n] \rightarrow T \quad \Gamma \vdash N_1 : S_1 \quad \Gamma \vdash N_2 : T_2 \quad \dots \quad \Gamma \vdash N_n : S_n}{\Gamma \vdash (M N_1 N_2 \dots N_n) : T} \text{T-APP} \\
\frac{\Gamma \vdash M_1 : T_1 \quad \Gamma \vdash M_2 : T_2 \quad \dots \quad \Gamma \vdash M_n : T_n}{\Gamma \vdash [M_1 M_2 \dots M_n] : [T_1 T_2 \dots T_n]} \text{T-TUPLE} \\
\frac{\Gamma \vdash M_1 : T \quad \Gamma \vdash M_2 : T \dots \Gamma \vdash M_n : T}{\Gamma \vdash \{M_1 M_2 \dots M_n\} : \{T\}} \text{T-COLLECTION} \\
\frac{\Gamma \vdash C : [S_1 S_2 \dots S_n] \rightarrow T \quad \Gamma \vdash N_1 : S_1 \quad \Gamma \vdash N_2 : S_2 \dots \Gamma \vdash N_n : S_n}{\Gamma \vdash \langle C N_1 N_2 \dots N_n \rangle : T} \text{T-INDUCTIVEDATA} \\
\frac{\Gamma \vdash M_1 : T_1 \quad \Gamma \vdash M_2 : (\text{Matcher } T_1) \quad \Gamma \vdash p : (\text{Pattern } T_1) \quad \Gamma \vdash V(\Gamma, p) \vdash M_3 : T_3}{\Gamma \vdash (\text{match-all } M_1 M_2 [p M_3]) : \{T_3\}} \text{T-MATCHALL} \\
\frac{}{\Gamma \vdash \text{something} : (\text{Matcher } T)} \text{T-SOMETHING} \\
\frac{\Gamma \vdash pp_i : (\text{PPPattern } T [S_k]_k) \quad \Gamma \vdash M_i : (\text{Matcher } [S_k]_k) \quad \Gamma \vdash dp_{ij} : (\text{PDPattern } T) \quad \Gamma \vdash V_{PP}(\Gamma, pp_i) \vdash V_{DP}(\Gamma, dp_i) \vdash N_{ij} : \{[S_k]_k\} \quad (\forall i, j)}{\Gamma \vdash (\text{matcher } [pp_i M_i [dp_{ij} N_{ij}]_j]_i) : (\text{Matcher } T)} \text{T-MATCHER}
\end{array}$$

Pattern:

$$\begin{array}{c}
\frac{}{\Gamma \vdash _ : (\text{Pattern } T)} \text{T-WILDCARD} \\
\frac{}{\Gamma \vdash \$x : (\text{Pattern } T)} \text{T-PATTERNVARIABLE} \\
\frac{\Gamma \vdash M : T}{\Gamma \vdash , M : (\text{Pattern } T)} \text{T-VALUEPATTERN} \\
\frac{\Gamma \vdash C_p : [(\text{Pattern } S_1) (\text{Pattern } S_2) \dots (\text{Pattern } S_n)] \rightarrow (\text{Pattern } T) \quad \Gamma \vdash M_1 : (\text{Pattern } S_1) \quad \Gamma \vdash V(\Gamma, M_1) \vdash M_2 : (\text{Pattern } S_2) \quad \Gamma \vdash V(\Gamma, M_1) \vdash V(\Gamma \vdash V(\Gamma, M_1), M_2) \vdash M_3 : (\text{Pattern } S_3) \quad \dots \quad \Gamma \vdash V(\Gamma, M_1) \vdash \dots \vdash V(\Gamma \vdash V(\Gamma, M_1) \vdash \dots, M_{n-1}) \vdash M_n : (\text{Pattern } S_n)}{\Gamma \vdash \langle C_p M_1 M_2 \dots M_n \rangle : (\text{Pattern } T)} \text{T-INDUCTIVEPATTERN}
\end{array}$$

Primitive pattern pattern:

$$\begin{array}{c}
\frac{}{\Gamma \vdash \$: (\text{PPPattern } T [T])} \text{T-PATTERNHOLE} \\
\frac{}{\Gamma \vdash , \$x : (\text{PPPattern } T [])} \text{T-VALUEPATTERNPATTERN} \\
\frac{\Gamma \vdash C_{pp} : [T_1 T_2 \dots T_n] \rightarrow T \quad \Gamma \vdash pp_i : (\text{PPPattern } T_i S_i)}{\Gamma \vdash \langle C_{pp} pp_i \rangle_i : (\text{PPPattern } T \Sigma_i S_i)} \text{T-INDUCTIVEPATTERNPATTERN}
\end{array}$$

Primitive data pattern:

$$\begin{array}{c}
\frac{}{\Gamma \vdash \$x : (\text{PDPattern } T)} \text{T-PRIMITIVEPATTERNVARIABLE} \\
\frac{\Gamma \vdash C_{dp} : [S_1 S_2 \dots S_n] \rightarrow T \quad \Gamma \vdash dp_1 : S_1 \quad \Gamma \vdash dp_2 : S_2 \dots \Gamma \vdash dp_n : S_n}{\Gamma \vdash \langle C_{dp} dp_1 dp_2 \dots dp_n \rangle : T} \text{T-INDUCTIVEDATAPATTERN}
\end{array}$$

V , V_{PP} , and V_{DP} return the type bindings for the variables in a pattern, primitive-pattern-pattern, and primitive-data-pattern,

respectively.

$$\begin{aligned}
V(\Gamma, _) &= \epsilon \\
V(\Gamma, \$x) &= \{(x : T)\} && \text{when } \Gamma \vdash \$x : (\text{Pattern } T) \\
V(\Gamma, , M) &= \epsilon \\
V(\Gamma, \langle C \ p_1 \ p_2 \ \dots \rangle) &= V(\Gamma, p_1) \ ++ \ V(\Gamma \ ++ \ V(\Gamma, p_1), p_2) \ ++ \ \dots
\end{aligned}$$

$$\begin{aligned}
V_{PP}(\Gamma, \$) &= \epsilon \\
V_{PP}(\Gamma, , \$x) &= \{(x : (\text{Pattern } T))\} && \text{when } \Gamma \vdash \$x : (\text{PPPattern } T) \\
V_{PP}(\Gamma, \langle C \ pp_1 \ pp_2 \ \dots \rangle) &= V_{PP}(\Gamma, pp_1) \ ++ \ V_{PP}(\Gamma, pp_2) \ ++ \ \dots
\end{aligned}$$

$$\begin{aligned}
V_{DP}(\Gamma, \$x) &= \{(x : T)\} && \text{when } \Gamma \vdash \$x : (\text{PDPattern } T) \\
V_{DP}(\Gamma, \langle C \ dp_1 \ dp_2 \ \dots \rangle) &= V_{DP}(\Gamma, dp_1) \ ++ \ V_{DP}(\Gamma, dp_2) \ ++ \ \dots
\end{aligned}$$

Q&A

- Does Egison type system support type polymorphism?
Let polymorphism has been implemented in our type system. But it is omitted to simplify this handout.

Examples

$$\begin{array}{c}
\frac{}{\epsilon \vdash \text{"Hello"} : \text{String}} \text{T-STR} \qquad \frac{}{\{x : \text{Integer}\} \vdash \text{"World!"} : \text{String}} \text{T-STR} \\
\\
\frac{\frac{}{\epsilon \vdash \text{true} : \text{Bool}} \text{T-BOOL} \quad \frac{}{\epsilon \vdash 10 : \text{Integer}} \text{T-NUM} \quad \frac{}{\epsilon \vdash 20 : \text{Integer}} \text{T-NUM}}{\epsilon \vdash (\text{if true } 10 \ 20) : \text{Integer}} \text{T-IF} \\
\\
\frac{\frac{\frac{}{\dots \vdash \text{b.} + : [\text{Integer Integer}] \rightarrow \text{Integer}} \text{T-VAR} \quad \frac{}{\dots \vdash x : \text{Integer}} \text{T-VAR} \quad \frac{}{\dots \vdash 10 : \text{Integer}} \text{T-NUM}}{\{\text{b.} + : [\text{Integer Integer}] \rightarrow \text{Integer}\} \ ++ \ \{x : \text{Integer}\} \vdash (\text{b.} + \ x \ 10) : \text{Integer}} \text{T-APP}}{\{\text{b.} + : [\text{Integer Integer}] \rightarrow \text{Integer}\} \vdash (\text{lambda } [\$x] (\text{b.} + \ x \ 10)) : \text{Integer} \rightarrow \text{Integer}} \text{T-ABS}
\end{array}$$